

Model categories of lcc categories and the gros model of dependent type theory

Martin E. Bidingmaier

Aarhus University

In this talk

An application of model category theory to the coherence problem of type theory.

1. Dependent type theory
 - ▶ Dependent type theory as essentially algebraic theory
 - ▶ Lcc categories and “gros” semantics
 - ▶ The coherence problem
2. Lcc sketches
 - ▶ Model categories of marked objects
 - ▶ Bousfield localization at “axioms”
3. Strict lcc categories
 - ▶ $sLcc$ as category of algebraically fibrant objects
 - ▶ Partial interpretation of type theory in $sLcc$
4. Algebraically cofibrant strict lcc categories
 - ▶ Strictification
 - ▶ A solution to the coherence problem
5. Conclusion & open problems

Dependent type theory

Extensional (!) dependent type theory is an essentially algebraic theory (CwFs).

Sorts:

- ▶ contexts Γ, Δ
- ▶ types σ, τ
- ▶ terms s, t
- ▶ morphisms f, g .

Operations:

- ▶ Every type σ is assigned its context: $\Gamma \vdash \sigma$.
- ▶ Every term s is assigned its context & type: $\Gamma \vdash s : \sigma$.
- ▶ Contexts and morphisms form a category.
- ▶ We can *substitute* terms and types along morphisms:

$$f : \Delta \rightarrow \Gamma \text{ and } \Delta \vdash s : \sigma \implies \Gamma \vdash f(s) : f(\sigma)$$

- ▶ Context extension:

$$p : \Gamma \rightarrow \Gamma.\sigma \qquad \Gamma.\sigma \vdash v : p(\sigma)$$

- ▶ Type formers:

$$\Gamma \vdash s_1, s_2 : \sigma \implies \Gamma \vdash \text{Eq } s_1 s_2 \qquad \Gamma.\sigma \vdash \tau \implies \Gamma \vdash \Pi_\sigma \tau$$

“Gros” semantics in lcc categories

Definition

A category \mathcal{C} is *locally cartesian closed (lcc)* if it has all finite limits and all pullback functors $f^* : \mathcal{C}_{/Y} \rightarrow \mathcal{C}_{/X}$ have adjoints $\Sigma_f \dashv f^* \dashv \Pi_f$.

Examples: Elementary toposes, ex/lex completions thereof.

If $f : X \rightarrow 1$ and $\phi \hookrightarrow X$:

$$\Pi_f(\phi) = \forall x : X, \phi(x)$$

$$\Sigma_f(\phi) \approx \exists x : X, \phi(x)$$

The category of all lcc categories should be model of type theory:

- ▶ Contexts Γ are lcc categories
- ▶ Types $\Gamma \vdash \sigma$ are objects $\sigma \in \text{Ob } \Gamma$
- ▶ Terms $\Gamma \vdash s : \sigma$ are morphisms $s : 1 \rightarrow \sigma$ in Γ .
- ▶ Morphisms $f : \Delta \rightarrow \Gamma$ are lcc functors.
- ▶ $\text{Eq } s_1 \ s_2$ is equalizer of $1 \begin{array}{c} \xrightarrow{s_1} \\ \xrightarrow{s_2} \end{array} \sigma$.
- ▶ $\Gamma.\sigma = \Gamma_{/\sigma}$
- ▶ $\Pi_\sigma \tau$ is given by $\tau \in \text{Ob } \Gamma_{/\sigma}$ and $\Pi_\sigma : \Gamma_{/\sigma} \rightarrow \Gamma$.

Coherence problems

- ▶ Substitution does not commute *strictly* with type formers:

$$f(\text{Eq } s_1 \ s_2) \cong \text{Eq } f(s_1) \ f(s_2)$$

- ▶ Context extension must have 1-categorical universal property, not bicategorical:

$$\begin{array}{ccc} & \Gamma/\sigma & \\ \sigma^* \nearrow & \cong & \searrow k^* \circ f/\sigma \\ \Gamma & \xrightarrow{f} & \Delta \end{array}$$

$$k : 1 \rightarrow f(\sigma) \text{ in } \Delta$$

Model categories

Definition

A *model category* is a category \mathcal{M} equipped with three classes of maps *cofibrations*, *fibrations*, *weak equivalences* subject to a number of axioms.

- ▶ Presents higher localization $\mathcal{W}^{-1}\mathcal{M}$ at weak equivalences \mathcal{W} .
- ▶ Cofibrations are “good monos”, fibrations “good epis”.
- ▶ Can always factor as $X \xrightarrow{\sim} X' \twoheadrightarrow Y$ and $X \twoheadrightarrow Y' \xrightarrow{\sim} Y$.
- ▶ In particular: $X \rightarrow 1$ (fibrant replacement) and $0 \rightarrow X$ (cofibrant replacement).
- ▶ *Combinatorial* model category: Locally presentable, generated by *sets* of (trivial) cofibrations.

Idea:

- ▶ Lcc categories form a higher category, presented as model category.
- ▶ Coherence problems are about underlying 1-category.
- ▶ Different (but Quillen equivalent) presentations as model categories can vary in underlying 1-categories.
→ find good presentation!

Marked objects

Definition (Isaev)

Let \mathcal{C} be a category and let $i : I \rightarrow \mathcal{C}$ be a diagram in \mathcal{C} . An (i) -marked object is given by an object $U(X)$ in \mathcal{C} and a set of morphisms of the form $k : i(K) \rightarrow U(X)$, the *marked* morphisms, such that

$$k : i(K_2) \rightarrow U(X) \text{ is marked and } f : K_1 \rightarrow K_2 \implies k \circ i(f) \text{ is marked.}$$

A morphism of i -marked objects is a marking-preserving morphism in \mathcal{C} .

- ▶ For us: $I \subseteq \mathcal{C} = \text{Cat}$ is subcategory, $\text{Cat}^i = \text{Cat}^I$.
- ▶ The forgetful functor $U : \mathcal{C}^I \rightarrow \mathcal{C}$ has both adjoints:

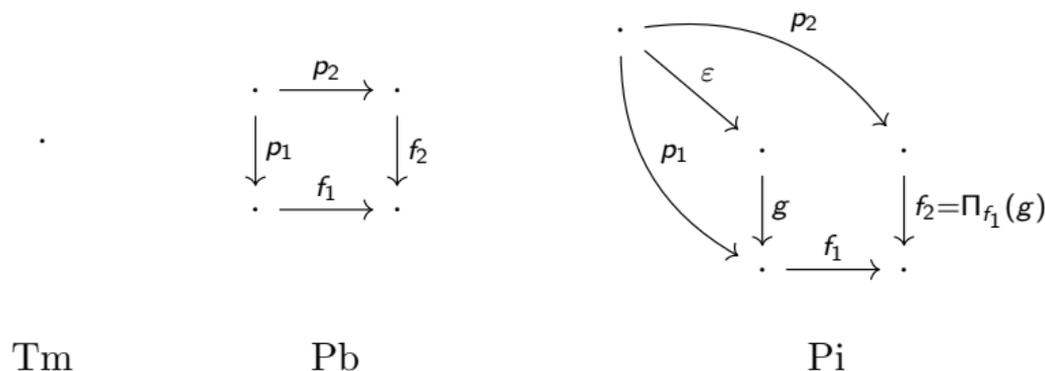
$$(-)^b \dashv U \dashv (-)^\sharp$$

Minimal^b and maximal[‡] markings.

Lcc shapes

Definition

$I_{\text{Lcc}} \subseteq \text{Cat}$ is given by the categories



and the inclusion $\text{Pb} \subseteq \text{Pi}$.

Intuition:

- ▶ Marked maps $\text{Tm} \rightarrow \mathcal{C}$: terminal objects,
- ▶ Marked maps $\text{Pb} \rightarrow \mathcal{C}$: pullback squares,
- ▶ Marked maps $\text{Pi} \rightarrow \mathcal{C}$: dependent products; $f_2 = \Pi_{f_1}(g)$.

Model category structure on marked objects

Theorem (Isaev)

Let \mathcal{M} be a combinatorial model category and let $i : I \rightarrow \mathcal{M}$ be a diagram in \mathcal{M} such that every object in the image of i is cofibrant. Then the following defines the structure of a combinatorial model category on \mathcal{M}^i :

- ▶ f in \mathcal{M}^i is a cofibration in \mathcal{M}^i iff $U(f)$ is cofibration in \mathcal{M} .
- ▶ f in \mathcal{M}^i is a weak equivalence iff $U(f)$ is a weak equivalence in \mathcal{M} and f reflects markings up to homotopy.

A marked object X is fibrant iff $U(X)$ is fibrant in \mathcal{M} and the markings of X are stable under homotopy. The adjunctions $(-)^b \dashv U$ and $U \dashv (-)^\sharp$ are Quillen adjunctions.

$\text{Cat}^{\text{lcc}} := \text{Cat}^{\text{Lcc}}$ inherits model structure from canonical model structure on Cat :

- ▶ Cofibrations are the functors that are injective on objects.
- ▶ Weak equivalences are marking-reflecting equivalences of categories.
- ▶ Fibrant objects are those where markings are stable under isomorphism of diagrams.

Lcc sketches

Idea: Restrict fibrant objects to those that are actual lcc categories.

→ add more trivial cofibrations.

Definition

The model category Lcc of *lcc sketches* is the left Bousfield localization $S^{-1}\text{Cat}^{lcc}$, where S consists of:



$$\left\{ \begin{array}{ccc} \cdot & \xrightarrow{p_2} & \cdot \\ \downarrow p_1 & & \downarrow f_2 \\ \cdot & \xrightarrow{f_1} & \cdot \end{array} \right\} \longrightarrow \left\{ \begin{array}{ccc} \cdot & \xrightarrow{p_2} & \cdot \\ \downarrow p_1 & \circlearrowleft & \downarrow f_2 \\ \cdot & \xrightarrow{f_1} & \cdot \end{array} \right\}$$

"Pullback squares commute."



$$\left\{ \begin{array}{ccc} & & \cdot \\ & & \downarrow f_2 \\ \cdot & \xrightarrow{f_1} & \cdot \end{array} \right\} \longrightarrow \left\{ \begin{array}{ccc} \cdot & \xrightarrow{p_2} & \cdot \\ \downarrow p_1 & & \downarrow f_2 \\ \cdot & \xrightarrow{f_1} & \cdot \end{array} \right\}$$

"All pullbacks exist."

Theorem

Lcc presents the higher category of lcc categories.

- ▶ *The fibrant objects are precisely lcc categories, with diagrams marked iff they satisfy universal property.*
- ▶ *Weak equivalences of fibrant objects are the equivalences of underlying categories.*
- ▶ *Ho Lcc is the category of lcc categories and isomorphism classes of lcc functors.*
- ▶ *The homotopy function complexes of fibrant lcc sketches are the groupoids of lcc functors and their isomorphisms.* □

The subcategory of fibrant lcc sketches is what's usually called "category of lcc categories".

Interpretation of type theory would depend on AC, suffers from coherence issues.

Question

Describe a set of generating (trivial) cofibrations for Lcc .

Algebraically fibrant objects

Definition

Let J be a suitable (e.g. generating) set of trivial cofibrations in a model category \mathcal{M} . The category $\text{Alg } \mathcal{M}$ of *algebraically fibrant objects* of \mathcal{M} (wrt. J) consists of object $G(X) \in \text{Ob } \mathcal{M}$ with assigned lifts

$$\begin{array}{ccc} A & \xrightarrow{a} & G(X) \\ \downarrow j & \nearrow \ell_X(j,a) & \\ B & & \end{array}$$

against all $j \in J$.

Model category structure on $\text{Alg } \mathcal{M}$

Theorem (Bourke)

Let \mathcal{M} be a combinatorial model category. Then $G : \text{Alg } \mathcal{M} \rightarrow \mathcal{M}$ has a left adjoint F . The model structure of \mathcal{M} can be transferred to $\text{Alg } \mathcal{M}$, and (F, G) is a Quillen equivalence. \square

- ▶ G reflects weak equivalences and fibrations.
- ▶ Every object in $\text{Alg } \mathcal{M}$ is fibrant.
- ▶ GF is fibrant replacement monad on \mathcal{M} .
- ▶ If every $X \in \text{Ob } \mathcal{M}$ is cofibrant: FG is cofibrant replacement comonad on $\text{Alg } \mathcal{M}$.
- ▶ Duality of property and structure.

Question

If \mathcal{M} is Gpd -enriched, then $\text{Alg } \mathcal{M}$ is Gpd -enriched. What about sSet -enrichment?

Question

What are the least requirements on J ? What if the lifts $\ell(j, a)$ are not specified for all a ?

Strict lcc categories

Definition

The category of *strict lcc categories* is given by

$$\mathbf{sLcc} = \mathbf{Alg Lcc}.$$

- ▶ Preservation of assigned lifts is trivial when lifts are unique
 \implies only choice of pullback, terminal objects, dependent products matter (but not maps induced by universal properties)
- ▶ Morphisms are *strict lcc functors*: Preserve lcc structure on the nose.
- ▶ *FG* is cofibrant replacement: Forget assigned lcc structure, freely adjoin new structure.

Question

Does \mathbf{sLcc} coincide with Lack's model category of algebras for a 2-monad T , instantiated with the free lcc category monad T on \mathbf{Cat} ?

Partial interpretation of type theory in $s\mathbb{L}cc$

- ▶ Contexts are $slcc$ categories Γ
- ▶ Morphisms are *strict* lcc functors $f : \Delta \rightarrow \Gamma$.
- ▶ Types $\Gamma \vdash \sigma$ are objects $\sigma \in \text{Ob } \Gamma$.
- ▶ Terms $\Gamma \vdash s : \sigma$ are morphisms $s : 1 \rightarrow \sigma$ in Γ .
- ▶ Finite limit types are defined by *canonical* finite limits in Γ .
- ▶ Strict substitution, e.g.

$$f(\text{Eq } s_1 \ s_2) = \text{Eq } f(s_1) \ f(s_2)$$

holds because morphisms in $s\mathbb{L}cc$ preserve lifts on the nose.

Context extensions vs slice categories

- ▶ Context extension $p : \Gamma \rightarrow \Gamma.\sigma$ is pushout

$$\begin{array}{ccc} F(\{t, \sigma\}) & \longrightarrow & F(\{v : t \rightarrow \sigma\}) \\ t \mapsto 1 \downarrow & & \downarrow \\ \Gamma & \xrightarrow{p} & \Gamma.\sigma \end{array}$$

- ▶ If $\tau \in \text{Ob } \Gamma.\sigma$, how to define $\Pi_\sigma \tau \in \text{Ob } \Gamma$?
- ▶ Want to apply $\Pi_\sigma : \Gamma_{/\sigma} \rightarrow \Gamma$.
- ▶ Do $\sigma^* : \Gamma \rightarrow \Gamma_{/\sigma}$ and diagonal $d : 1 = \text{id}_\sigma \rightarrow \sigma^*(\sigma)$ induce map $\Gamma.\sigma \rightarrow \Gamma_{/\sigma}$?
- ▶ No, σ^* is map in Lcc , not in sLcc (not strict).

Strictification

$\Gamma \in \text{ObsLcc}$ is cofibrant $\iff \varepsilon : F(G(\Gamma)) \rightarrow \Gamma$ is retraction:

$$\exists \lambda : \Gamma \rightarrow F(G(\Gamma)), \varepsilon \lambda = \text{id}$$

Now:

$$\sigma^* : G(\Gamma) \rightarrow G(\Gamma/\sigma)$$

$$\rightsquigarrow \overline{\sigma^*} : F(G(\Gamma)) \rightarrow \Gamma/\sigma$$

$$\rightsquigarrow (\sigma^*)^s := \overline{\sigma^*} \circ \lambda : \Gamma \rightarrow F(G(\Gamma)) \rightarrow \Gamma/\sigma$$

Proposition

If $\Gamma \in \text{sLcc}$ is cofibrant, then for all $f : G(\Gamma) \rightarrow G(\Delta)$ there exists $f^s : \Gamma \rightarrow \Delta$ such that $f \cong G(f^s)$. □

Good: Have map $a : \Gamma.\sigma \rightarrow \Gamma/\sigma$.

Bad: a is not compatible with morphisms in Γ ! Need

$$g : E \rightarrow \Gamma \implies f^s \circ g = (f \circ G(g))^s$$

Algebraically cofibrant objects

Definition

Let \mathcal{M} be a model category and let C be a cofibrant replacement comonad on \mathcal{M} . An *algebraically cofibrant object* of \mathcal{M} is a coalgebra for C ; the category of such objects is denoted by $\text{Coa } \mathcal{M}$.

- ▶ Structure map $\lambda : X \rightarrow C(X)$ is inclusion of retract
 \implies coalgebras are cofibrant in sLcc .
- ▶ Coalgebra morphisms preserve λ .

Theorem (Ching & Riehl)

Let \mathcal{M} be a combinatorial and simplicial model category. Then there exists a suitable simplicial cofibrant replacement comonad. The model category structure of \mathcal{M} can be transferred to $\text{Coa } \mathcal{M}$ such that \mathcal{M} and $\text{Coa } \mathcal{M}$ are Quillen equivalent.

The interpretation of type theory in Coa sLcc

- ▶ Category of contexts is Coa sLcc .
- ▶ Types, terms, finite limit types are interpreted as in sLcc .
- ▶ $\text{Coa sLcc} \rightarrow \text{sLcc}$ commutes with context extension.

Lemma

There is a natural transformation of functors $(\text{Coa sLcc})_ \rightarrow \text{sLcc}$*

$$((\lambda, \Gamma, \sigma) \mapsto \Gamma.\sigma) \Rightarrow ((\lambda, \Gamma, \sigma) \mapsto \Gamma/\sigma).$$

whose components are weak equivalences in sLcc

- ▶ Homotopy inverses can be found in Lcc , also strictly natural.

Theorem

The opposite of Coa sLcc carries cwf structure (i.e. is a model of type theory) that supports finite limit, Π and Σ types.

Coherence problems are about the underlying 1-category of model categories.
Quillen equivalent model categories can vary in underlying 1-categories.

1. *Model category of sketches:*

Universal objects *merely exist*, no canonical choice.

→ cannot even *state* substitution stability.

2. *Algebraically fibrant objects:*

Have canonical universal objects preserved by morphisms/substitution.

But: Context extension $\Gamma.\sigma$ is pushout, only “correct” when Γ is cofibrant.

3. *Algebraically cofibrant objects:*

Cofibrancy baked into structure.

Can strictify maps $f : G(\Gamma) \rightarrow G(\Delta)$, strictification functorial.

Some open problems

Question

Can some very weak variant of HoTT be interpreted in lcc quasi-categories with this technique?

Question

Is there a model category of sketches for every 2-monad T on \mathbf{Cat} ? What about \mathbf{sSet} -enriched monads on \mathbf{sSet} ?

Marked objects probably only work when T -algebra structure is essentially unique ($\rightarrow T$ is modality), but not e.g. for monoidal categories.

Question

Quasi-categories/Kan complexes are weird: Composition is property and preserved up to homotopy, identities are structure and preserved up to equality. Is \mathbf{sSet} of the form $\mathbf{Alg}_{\mathbf{sSet}}$ (or variation) for a model category of semi-simplicial sets \mathbf{ssSet} ?

References

- [Bid20] Martin E. Bidlingmaier, *An interpretation of dependent type theory in a model category of locally cartesian closed categories*, 2020.
- [Bou19] John Bourke, *Equipping weak equivalences with algebraic structure*, *Mathematische Zeitschrift* (2019).
- [CR14] Michael Ching and Emily Riehl, *Coalgebraic models for combinatorial model categories*, *Homology, Homotopy and Applications* **16** (2014), no. 2, 171–184.
- [Isa16] Valery Isaev, *Model category of marked objects*, 2016.
- [Lac07] Stephen Lack, *Homotopy-theoretic aspects of 2-monads.*, *Journal of Homotopy and Related Structures* **2** (2007), no. 2, 229–260.
- [Vic16] Steven Vickers, *Sketches for arithmetic universes*, *Journal of Logic and Analysis* (2016).